

1. (15 %) Suppose a thread is running in a critical section of code. It means that the thread has acquired all the locks through proper arbitration. Can this thread get context switched? Please explain the reasons.
2. (20 %) Suppose that a disk drive has 5000 cylinders, numbered 0 to 4999. The drive is currently serving a request at cylinder 134, and the previous request was at cylinder 125. The queue of pending requests, in FIFO order, is 86, 1470, 913, 3774, 948, 4509, 1022, 2750, 130. Starting from the current head position, what is the total distance (in cylinders) that the disk arm moves to satisfy all the pending requests, for SSTF and SCAN disk scheduling algorithms?
3. (15 %) Are the following statements about IP addresses true or false? For each statement, you will get 3 points for correct answer, zero point for blank, or -2 point for incorrect answer.
  - (a) The subnet mask for the subnet 200.23.16.0/23 is 255.255.255.0.
  - (b) The subnet 200.23.16.0/23 could accommodate up to 256 hosts.
  - (c) Domain Name Service (DNS) can be used to acquire IP addresses.
  - (d) Address Resolution Protocol (ARP) can be used to acquire IP addresses.
  - (e) Network Address Translation (NAT) is used to map MAC addresses to IP addresses.
4. (30%) In the consumer-producer example program, a ring buffer queue is used to store the produced item that will be taken off by the consumer later.  
There are two example programs, the following one (figure 5.9, 5.10) uses semaphore to implement it. (p240, Fig 6.10/6.11 8<sup>th</sup> edition)  
4-a.(5%) What are the shared variables in figure 5.9, 5.10?  
4-b.(5%) Why this program needs to use semaphore primitive ?

```
int n;
semaphore mutex = 1;
semaphore empty = n;
semaphore full = 0

do {
    . . .
    /* produce an item in next_produced */
    . . .
    wait(empty);
    wait(mutex);
    . . .
    /* add next_produced to the buffer */
    . . .
    signal(mutex);
    signal(full);
} while (true);
```

Figure 5.9 The structure of the producer process.

```

do {
    wait(full);
    wait(mutex);
    . . .
    /* remove an item from buffer to next_consumed */
    . . .
    signal(mutex);
    signal(empty);
    . . .
    /* consume the item in next_consumed */
    . . .
} while (true);

```

Figure 5.10 The structure of the consumer process.

4-c.(5%) The following figure 3.13 and 3.14 are another example program. What are the shared variables in these producer and consumer program?

4-d.(5%) In what conditions the program need NOT to use the lock synchronization primitive to support the correct processing? (p118, fig 3.14/3.15 8<sup>th</sup> edition)

```

#define BUFFER_SIZE 10

typedef struct {
    . . .
} item;

item buffer[BUFFER_SIZE];
int in = 0;
int out = 0;

while (true) {
    /* produce an item in next_produced */

    while (((in + 1) % BUFFER_SIZE) == out)
        ; /* do nothing */

    buffer[in] = next_produced;
    in = (in + 1) % BUFFER_SIZE;
}

```

Figure 3.13 The producer process using shared memory.

```
item next_consumed;

while (true) {
    while (in == out)
        ; /* do nothing */

    next_consumed = buffer[out];
    out = (out + 1) % BUFFER_SIZE;

    /* consume the item in next_consumed */
}
```

Figure 3.14 The consumer process using shared memory.

4-e.(5%) If kernel monitor(R.C.A. Hoare Monitor) approach is adopted to support producer and consumer program. Any processes that can not call consumer or producer function in monitor simultaneously. That means the monitor restricts the parallel processing of these two functions. What is your suggestion to support multiprocessor parallel processing?

4-f.(5%). Why the problems addressed in the chapter Synchronization are nonexistent in functional languages, but there are existed in any procedural languages.

5. (20%) Show how to implement the wait() and signal() semaphore operations in multiprocessor environments using the test\_and\_set() instruction. The solution should exhibit minimal busy waiting.