

1. A prune-and-search algorithm consists of several iterations. At each iteration, it prunes away a fraction  $f$ ,  $0 < f < 1$ , of input data, and then it invokes the same algorithm recursively to solve the problem for the remaining data. After a certain number of iterations, the size of input data will be so small that the problem can be solved directly with constant time. Assume that the time complexity of a prune-and-search algorithm is  $T(n)$  and the time needed to execute each iteration is  $O(n^k)$  for some constant  $k$ ,  $k > 0$ . Show that  $T(n) = O(n^k)$ . (20%)
2. Show that the average case lower bound of the sorting problem is  $\Omega(n \log n)$ . (15%)
3. Given a problem  $X$ , how can we prove it to be NP-hard? (15%)
4. Given a list of  $n$  positive integers, you are asked to partition the list into two sublists, each of size  $\lfloor n/2 \rfloor$  or  $\lceil n/2 \rceil$ , such that the difference between the sums of the integers in the two sublists is minimized. (a) Give a decision version of this problem, and show that it is NP-complete. (15%) (b) If the summation of these  $n$  positive integers is small enough, it is possible to solve this problem efficiently. Design such an algorithm. (15%)
5. For two points  $P = (p_1, p_2)$ , and  $Q = (q_1, q_2)$  in the plane, we say that  $P$  **dominates**  $Q$  if  $p_1 > q_1$  and  $p_2 > q_2$ . Given a set  $S$  of  $n$  points, the **rank** of a point  $P$  in  $S$  is the number of points in  $S$  dominated by  $P$ . The problem is to find the rank of every point in  $S$ . A straightforward way to solve this problem is to conduct an exhaustive comparison of all pairs of points. Hence, this approach requires  $O(n^2)$  running time. Design a faster algorithm (i.e. the running time of your algorithm must be in  $o(n^2)$  order) to solve this problem. (20%)