國立中央大學資訊工程學系博士班 101 學年度第一學期資格考試題紙

## 科目: 作業系統 (Operating System) 第一頁 共二頁 (pagel of 2)

1. (21%) Consider the following snapshot of a system:

	Allocation	<u>Matrix</u>	<u>Available</u>
	ABCD	ABCD	ABCD
P0	0012	0012	1520
P1	$1\ 0\ 0\ 0$	1750	
P2	1354	2356	
P3	0632	0652	
P4	0014	0656	

Answer the following questions using the banker's algorithm. Note that all the following questions are independent.

- (a) Is the system in a safe state? Be sure to justify your answer.
- (b) If a request from process P1 arrives for (0, 4, 2, 0), can the request be granted immediately? <u>Be sure to justify your answer.</u>
- (c) If a request from process P4 arrives for (0, 4, 2, 0), can the request be granted immediately? <u>Be sure to justify your answer.</u>
- 2. (20%) For the dining philosophers problem, assume it has reached the deadlock state that each philosopher,  $P_i$ , has acquired the chopstick on his right-hand side,  $C_i$ , where  $1 \le i \le 5$ . Please draw (a) the resource allocation graph and (b) the wait-for graph for the philosopher deadlock situation.



- 3. (9%) What are the advantages and disadvantages of using circuit switching?
- 4. (10%) What are two differences between user-level threads and kernel-level threads? Under what circumstances is one type better than the other?
- 5. (15%)Explain the purpose of the checkpoint mechanism. 5%

How often should checkpoints be performed? Describe how the frequency of checkpoints affects: 10%

- System performance when no failure occurs
- The time it takes to recover from a system crash
- The time it takes to recover from a disk crash

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## 科目: 作業系統 (Operating System) 第二頁 共二頁(page 2 of 2)

6. (25%)The first known correct software solution to the critical-section problem for n processes with a lower bound on waiting of n-1 turns was presented by Eisenberg and McGuire. The Processes share the following variables: 1. enum pstate { idle, want\_in, in\_cs}; 2. pstate flag[n]; 3. int turn; 4. 5. All the elements of flag are initially idle; the initial value of turn is immaterial(between 0 and n-1). The structure of process Pi is shown in the following: 6. do { 7. while(TRUE) { 8. flag[i] = want\_in; 9. i = turn;10. while (j != i) { 11. if (flag[j] != idle) { j = turn;12. 13. else 14. j = (j+1) % n;15. } 16. flag[i] = in\_cs; 17. i = 0;18. while  $((j < n) \&\& (j == i \parallel flag[j] != in_cs)) j++;$ 19. if  $((i \ge n) \&\& (turn == i || flag[turn] == idle))$  break; 20. } 21. 22. turn = i;23. // CRICIAL-SECTION 24. 25. 26. j = (turn+1) % n;27. while (flag[j] == idle) 28. j = (j+1) % n;29. turn = j; 30. flag[i] = idle; 31. 32. // REMAINDER-SECTION 33. } while(TRUE); 34.

(1).Prove that the algorithm satisfies all three requirements for the critical section. 15% (2).What is the basic assumption for hardware requirements? If the system is multicore or multiprocessor system, the processor use the cache to speed up the reference of shared memory , what problems should be considered ? 10%