

(12) (7 points) Write EBNF and syntax graph descriptions for the following:

(i). A Java class definition header statement

(ii). A C switch statement

Ans.

(i) $\langle \text{class_head} \rangle \rightarrow \{ \langle \text{modifier} \rangle \} \text{ class } \langle \text{id} \rangle [\text{extends class_name}]$
 $[\text{implements } \langle \text{interface_name} \rangle \{, \langle \text{interface_name} \rangle \}]$

$\langle \text{modifier} \rangle \rightarrow \text{public} \mid \text{abstract} \mid \text{final}$

(ii) $\langle \text{switch_stmt} \rangle \rightarrow \text{switch} (\langle \text{expr} \rangle) \{ \text{case } \langle \text{literal} \rangle : \langle \text{stmt_list} \rangle$
 $\{ \text{case } \langle \text{literal} \rangle : \langle \text{stmt_list} \rangle \} [\text{default} : \langle \text{stmt_list} \rangle] \}$

(13) (7 points) Rewrite the BNF of Example 3.4 to give + precedence over * and force + to be right associative.

Ans.

$\langle \text{assign} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle$

$\langle \text{id} \rangle \rightarrow A \mid B \mid C$

$\langle \text{expr} \rangle \rightarrow \langle \text{expr} \rangle * \langle \text{term} \rangle$
 $\mid \langle \text{term} \rangle$

$\langle \text{term} \rangle \rightarrow \langle \text{factor} \rangle + \langle \text{term} \rangle$
 $\mid \langle \text{factor} \rangle$

$\langle \text{factor} \rangle \rightarrow (\langle \text{expr} \rangle)$
 $\mid \langle \text{id} \rangle$

(14) (7 points) Modify the grammar of Example 3.4 to add a unary minus operator that has higher precedence than either + or *.

Ans.

Assume that the unary operators can precede any operand.

Replace the rule

$\langle \text{factor} \rangle \rightarrow \langle \text{id} \rangle$

with

$\langle \text{factor} \rangle \rightarrow \langle \text{id} \rangle$
 $\mid - \langle \text{id} \rangle$

(15) (7 points) Write an attribute grammar whose BNF basis is that of Example 3.6 in Section 3.4.5, but whose language rules are as follows: Data types cannot be mixed in expressions, but assignment statements need not have the same types on both sides of the assignment operator.

Ans.

1. Syntax rule: $\langle \text{assign} \rangle \rightarrow \langle \text{var} \rangle = \langle \text{expr} \rangle$

2. Syntax rule: $\langle \text{expr} \rangle \rightarrow \langle \text{var} \rangle[2] + \langle \text{var} \rangle[3]$

predicate: $\langle \text{var} \rangle[2].\text{actual_type} = \langle \text{var} \rangle[3].\text{actual_type}$

3. Syntax rule: $\langle \text{expr} \rangle \rightarrow \langle \text{var} \rangle$

4. Syntax rule: $\langle \text{var} \rangle \rightarrow \mathbf{A} \mid \mathbf{B} \mid \mathbf{C}$

Semantic rule; $\langle \text{var} \rangle.\text{actual_type} \leftarrow \text{look-up}(\langle \text{var} \rangle.\text{string})$